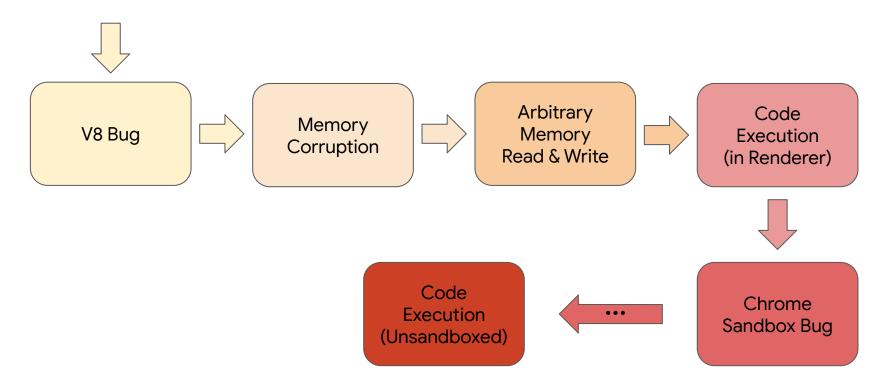
The V8 Heap Sandbox

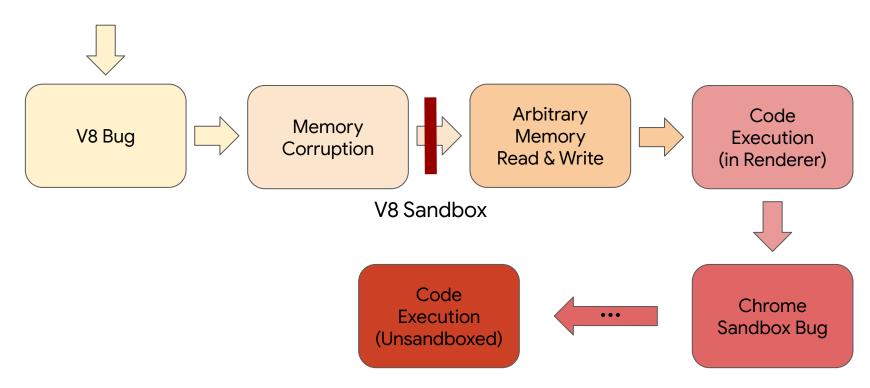
OffensiveCon 2024

Samuel Groß - Google V8 Security

Typical Exploit Flow

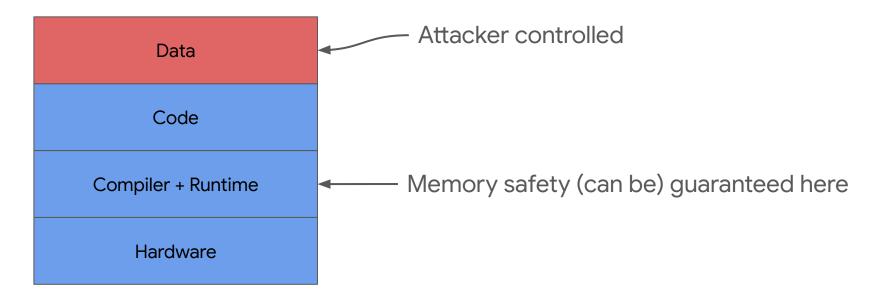


Typical Exploit Flow



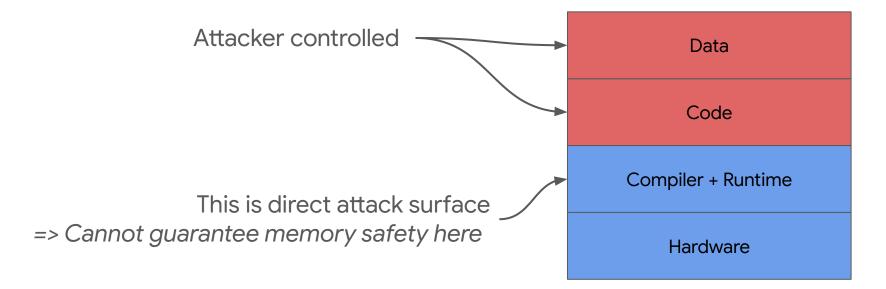
Why JavaScript Engine Security is hard

"Typical" Application



Why JavaScript Engine Security is hard





Why JavaScript Engine Security is hard

- Compiler-based memory safety doesn't work if compiler is attack surface
 - => Any logic bug can potentially turn into memory corruption
- Disabling optional compilers solves only a part of the problem
 - Plenty of bugs elsewhere (e.g. runtime) ...
 - ... and it is very slow :(

=> Writing a high-performance, memory-safe JS engine is **hard**

High-performance, memory-safe JavaScript engine?

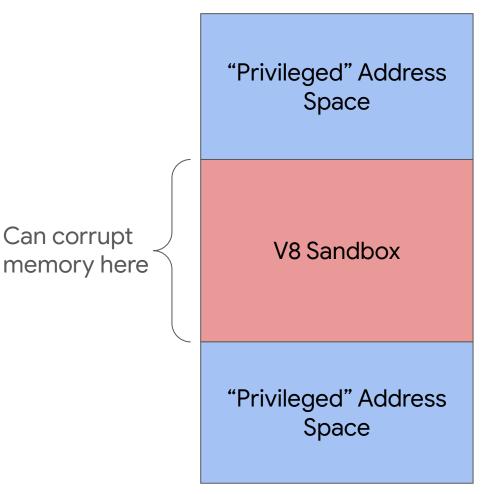
Write Bug Free Code

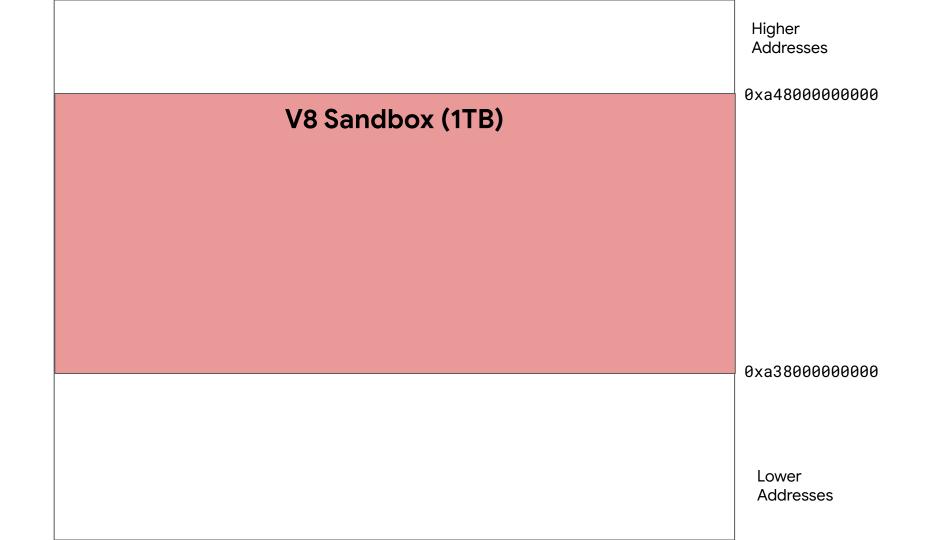
Big, hard problem ---

A different approach...

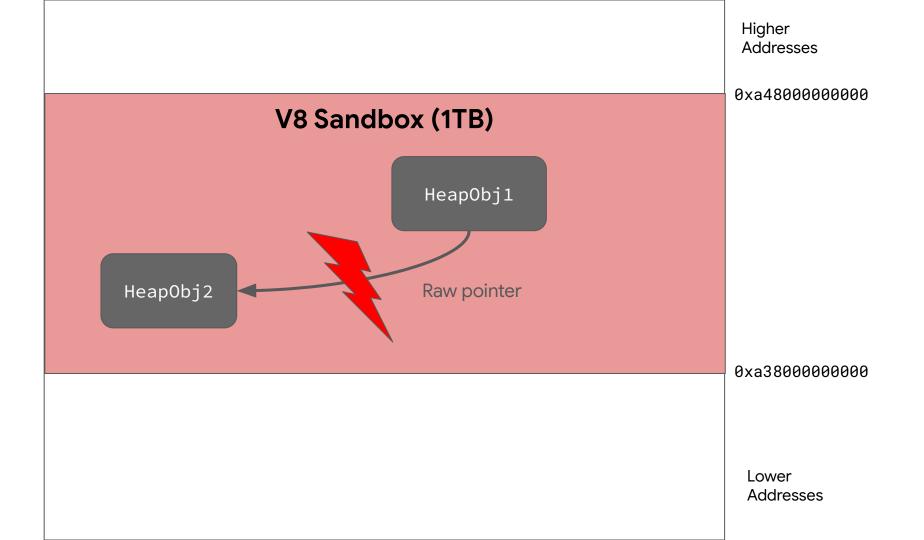
Idea:

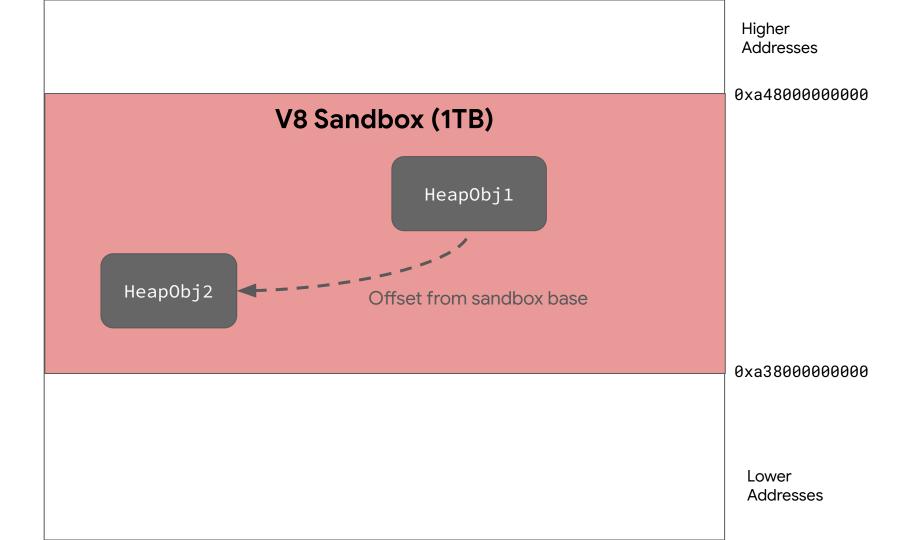
- Accept that bugs will happen and that memory will be corrupted
- Limit which memory can be corrupted
- Make that a security boundary
- => Result: an in-process sandbox

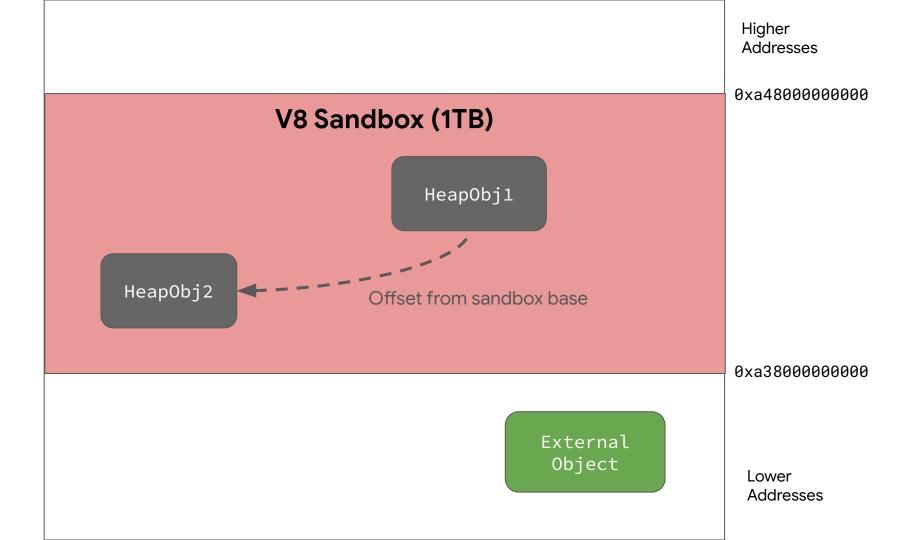


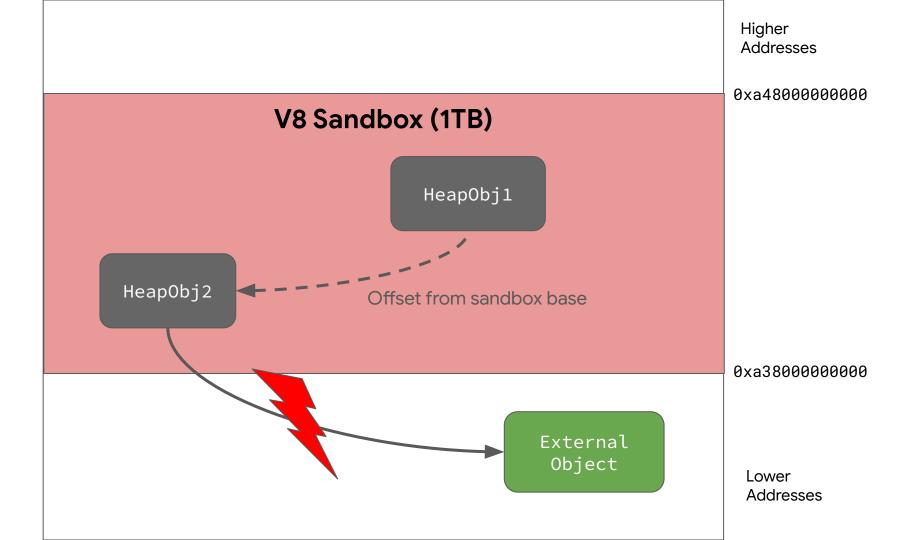


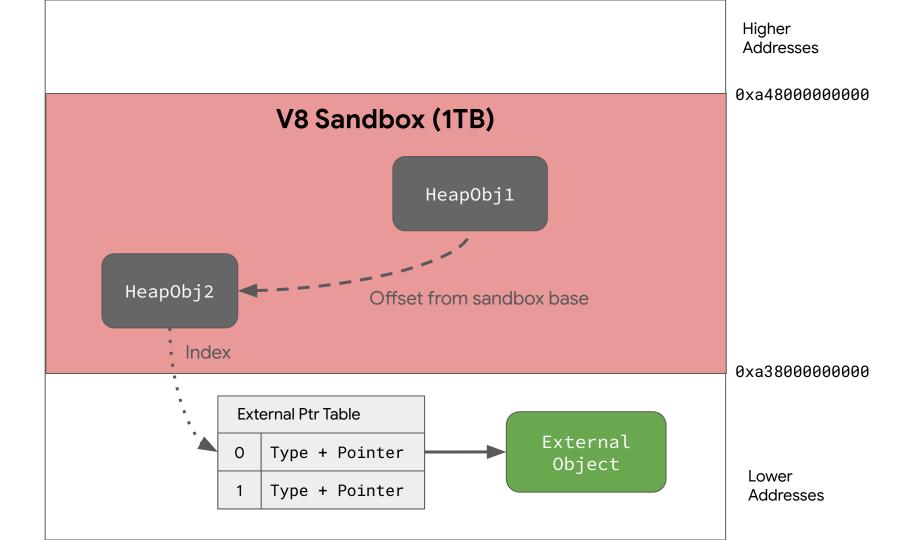








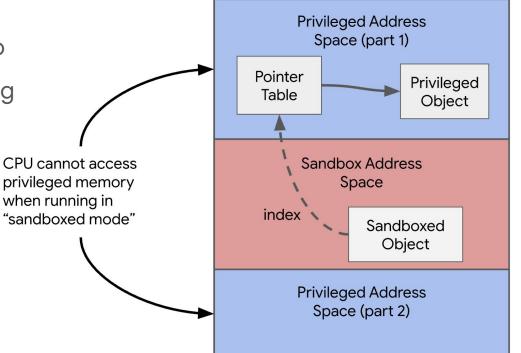






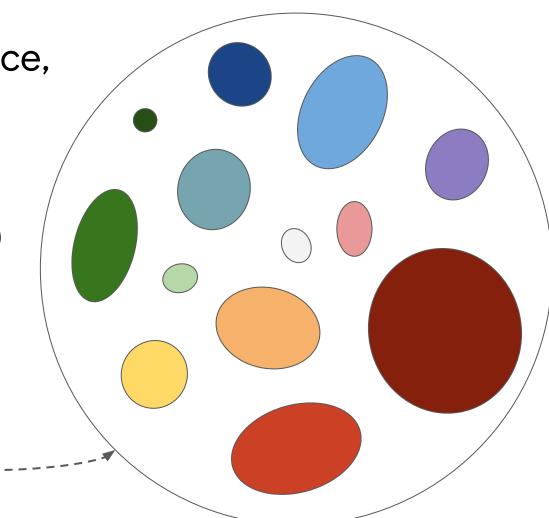
Sandbox with Hardware Support?

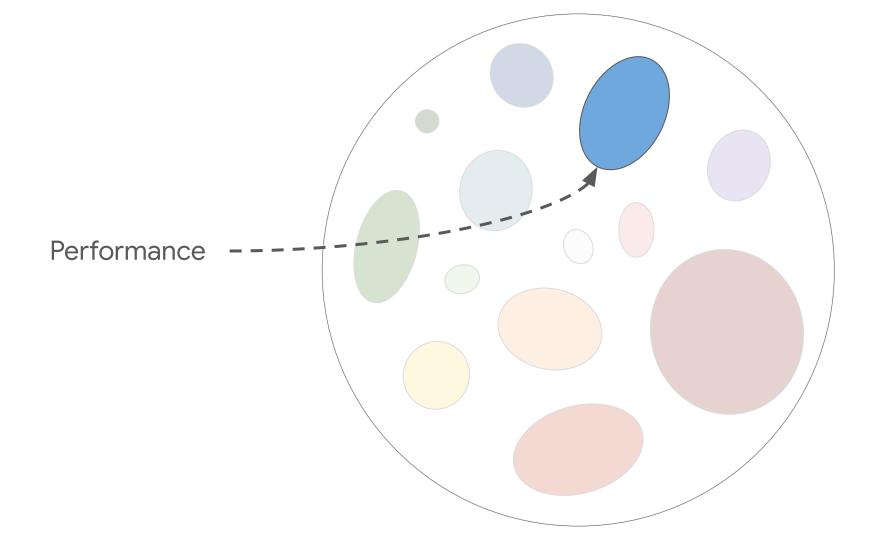
- In the future, may be possible to "drop privileges" when executing JS or Wasm code
- Would be very similar to userspace/kernel split
- Ideally: want to be able to run untrusted *machine code*



High-performance, memory-safe JavaScript engine! (with a sandbox)

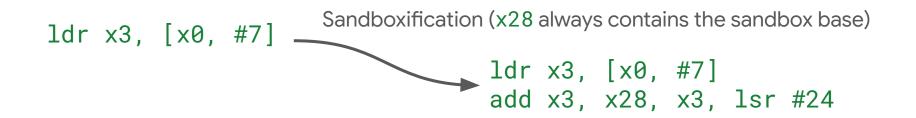
> Lots of smaller, simpler problems

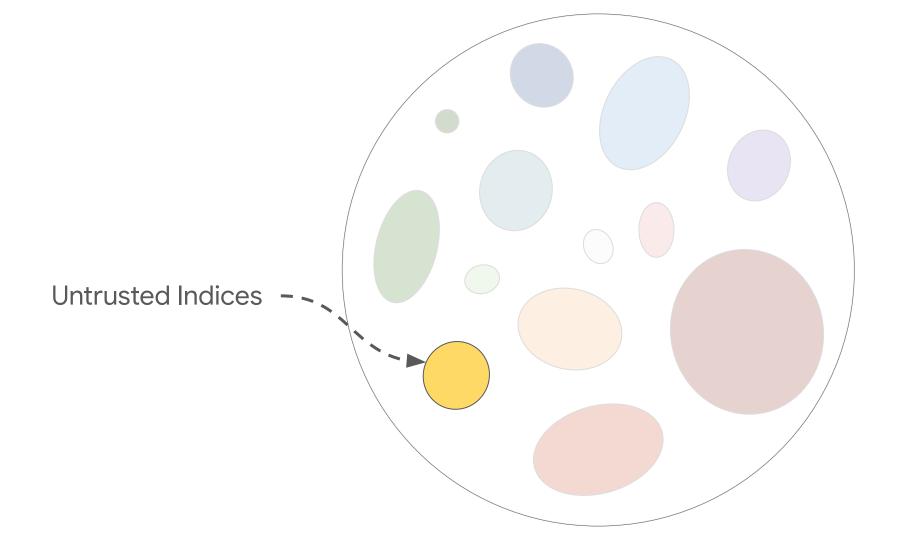




Performance

- Sandbox building blocks are fundamentally cheap
 - Offsets require just an additional add or shift+add instruction
 - Pointer table requires one additional memory load for external references
- => Benefit over other memory safety technologies
- Today: overhead of sandbox is only around 1% on popular benchmarks
 - => Can be (and is already) enabled by default!



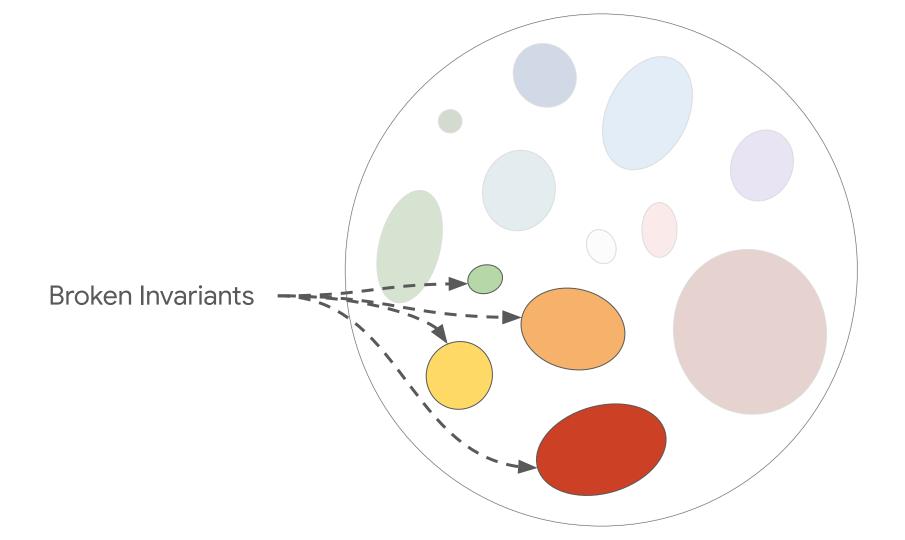


Untrusted Indices

Tagged<MyHeapObject> obj = ...; int idx = obj->get_the_index(); int val = obj->get_the_value(); some_global_array[idx] = val;

Untrusted Indices

Tagged<MyHeapObject> obj = ...; uint idx = obj->get_the_index(); int val = obj->get_the_value(); SBXCHECK(idx < some_global_array_size);</pre> some_global_array[idx] = val;



Broken Invariants

std::vector<std::string> JSObject::GetPropertyNames() { int num_properties = TotalNumberOfProperties(); std::vector<std::string> properties(num_properties);

```
for (int i = 0; i < NumberOfInObjectProperties(); i++) {
    properties[i] = GetNameOfInObjectProperty(i);
}</pre>
```

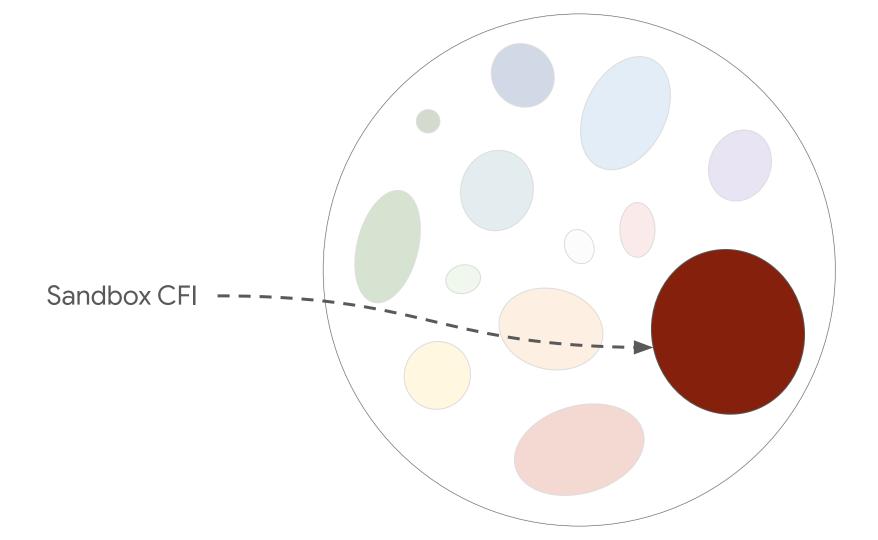
```
// Deal with the other types of properties
// ...
```

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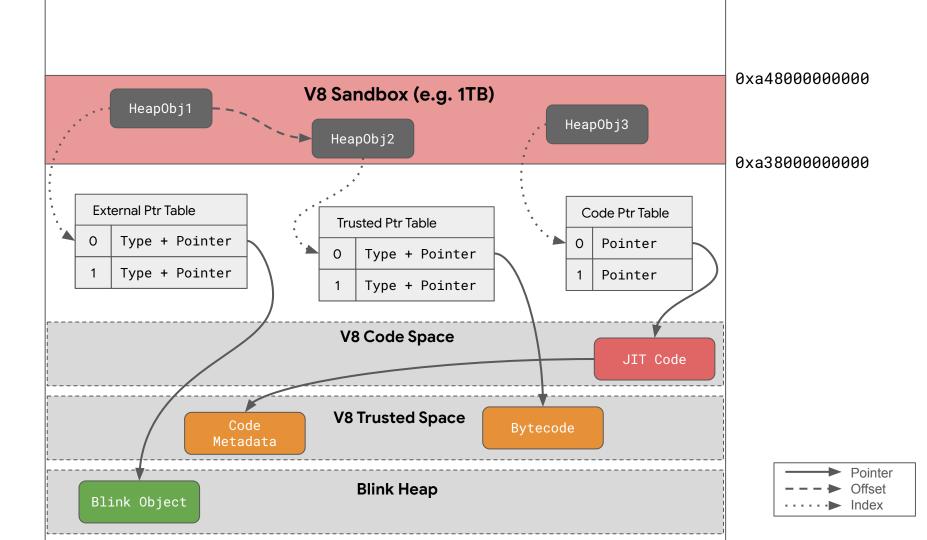


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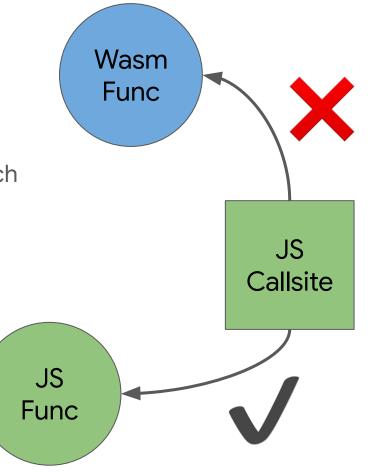
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 - Can e.g. lead to code corruption when manipulated
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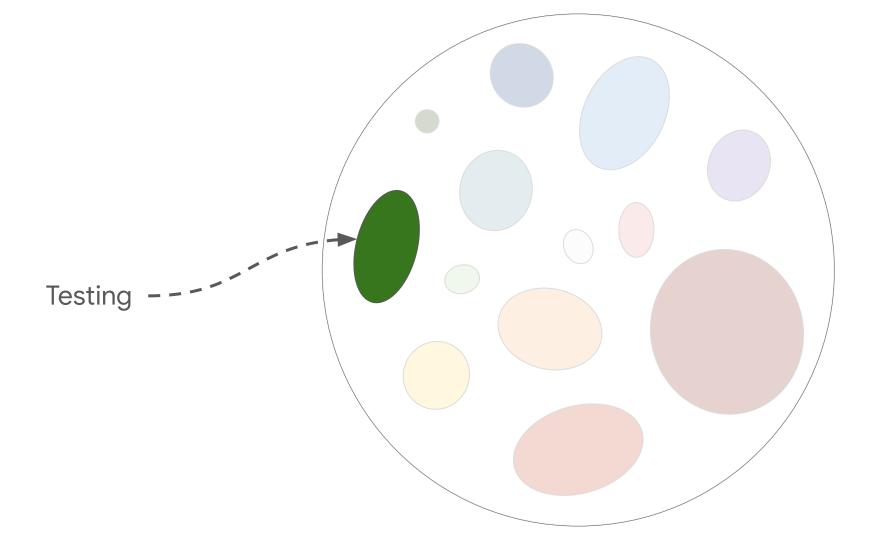
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 - Can e.g. lead to code corruption when manipulated
 - => Move out of the sandbox
- Less obvious: interpreter bytecode cannot be in the sandbox
 - Causes stack corruption if manipulated
 - => Move out of sandbox and also reference via a pointer table



And more subtle issues in this area:

- Calling convention/signature mismatch
- Deoptimization and tier-up
- Desynchronized code references
- ...
- => Still work to do in this area





Testing

- Sandbox is testable
 - Clear attacker model + tools to develop and validate sandbox bypasses
- This enables:
 - automatic fuzzing
 - ability to write regression tests
 - inclusion in Chrome's bug bounty program (active since March 2024)

let memory = new Sandbox.MemoryView(0, kSize); let dv = new DataView(memory); // Full read+write to sandbox address space dv.setUint8(0x41414141, 0x42);

Demo

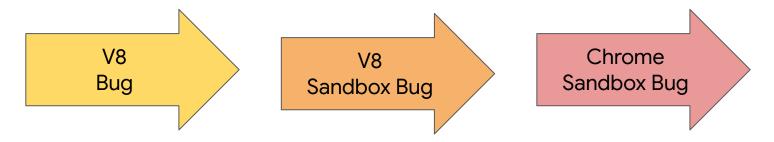
```
saelo v8/v8 [main] > cat demo.js
const kSize = 0x100000000; // 4 GB
let memory = new DataView(new Sandbox.MemoryView(0, kSize));
memory.setUint32(0x41414141, 0x42424242);
```

saelo v8/v8 [main] > ./out/x64.sbxtst/d8 --s



Conclusion

Sandbox increases length of (typical) V8-based Chrome exploit chain



Key question: how hard is this new attack surface?

... Only one way to find out: build it, then see what happens :)

Resources

- Blog post: <u>v8.dev/blog/sandbox</u>
- README: src/sandbox/README.md
- Past sandbox bugs: <u>v8-sandbox buganizer hotlist</u>
- Sandbox VRP rules: <u>g.co/chrome/vrp/#v8-sandbox-bypass-rewards</u>

Questions?